

Solutions - Midterm Exam

(February 15th @ 7:30 pm)

Presentation and clarity are very important! Show your procedure!

PROBLEM 1 (19 PTS)

- Compute the result of the following operations. The operands are signed fixed-point numbers. The result must be a signed fixed-point number. For the division, use $x = 5$ fractional bits.

$\begin{array}{r} 1.0111 + \\ 1.101010 \\ \hline \end{array}$	$\begin{array}{r} 1.010110 - \\ 1000.0101 \\ \hline \end{array}$	$\begin{array}{r} 01.11111 + \\ 0.10011 \\ \hline \end{array}$
$\begin{array}{r} 1.01101 \times \\ 10.101 \\ \hline \end{array}$	$\begin{array}{r} 0.111 \times \\ 1.0101 \\ \hline \end{array}$	$\begin{array}{r} 01.011 \div \\ 1.011 \\ \hline \end{array}$

$$\begin{array}{c} \text{C}_8 \text{C}_7 \text{C}_6 \text{C}_5 \text{C}_4 \text{C}_3 \text{C}_2 \text{C}_1 \text{C}_0 \\ 1.0111 + \\ 1.101010 \\ \hline 1.000110 \end{array} \quad \begin{array}{c} 1111.0101 - \\ 1000.0101 \\ \hline 1.000000 \end{array} \Rightarrow \begin{array}{c} \text{C}_8 \text{C}_7 \text{C}_6 \text{C}_5 \text{C}_4 \text{C}_3 \text{C}_2 \text{C}_1 \text{C}_0 \\ 1111.0101 + \\ 0111.1010 \\ \hline 0111.0000 \end{array} \quad \begin{array}{c} \text{C}_8 \text{C}_7 \text{C}_6 \text{C}_5 \text{C}_4 \text{C}_3 \text{C}_2 \text{C}_1 \text{C}_0 \\ 001.1111 + \\ 000.1001 \\ \hline 010.1010 \end{array}$$

$$\begin{array}{r} 10.101 \times \\ 1.01101 \\ \hline 011010001 \end{array} \Rightarrow \begin{array}{r} 01.011 \times \\ 0.10011 \\ \hline 011010001 \end{array} \Rightarrow \begin{array}{r} 10011 \times \\ 1011 \\ \hline 10011 \\ 10011 \\ 00000 \\ \hline 10011 \\ 011010001 \\ \hline 0.11010001 \end{array} \quad \begin{array}{r} 0.111 \times \\ 1.0101 \\ \hline 0.111 \\ 0.111 \\ \hline 1011 \\ 1011 \\ \hline 01001101 \\ \hline 0.1001101 \\ \hline 1.0110011 \end{array}$$

- ✓ $\frac{01.011}{1.011}$: To unsigned (denominator) and then alignment, $a = 3$: $\frac{01.011}{0.101} = \frac{01.011}{00.101} \equiv \frac{1011}{101}$

$$\begin{array}{r} 001000110 \\ 101 \overline{) 101100000} \\ \underline{101} \\ 01000 \\ \underline{101} \\ 110 \\ \underline{101} \\ 10 \end{array}$$

Append $x = 5$ zeros: $\frac{101100000}{101}$

Integer Division:

$$Q = 1000110, R = 10 \\ \rightarrow Qf = 10.00110 (x = 5)$$

$$\text{Final result (2C): } \frac{10.1010}{0.101} = 2C(010.0011) = 101.1101$$

PROBLEM 2 (11 PTS)

- Represent these numbers in Fixed Point Arithmetic (signed numbers). Use the FX format [8 4]. Truncate (the LSB) and perform Saturation when required.

✓ -33.375

$$+33.375 = 0100001.011 \Rightarrow -33.375 = 1011110.101$$

$$-33.375 = 1011110.101 = 1000.0000 \text{ (saturation)}$$

✓ 17.875

$$+17.875 = 010001.111$$

$$+17.875 = 0111.1111 \text{ (sat+truncation)}$$

- Complete the table for the following fixed-point formats (signed numbers): (3 pts.)

Integer bits	Fractional Bits	FX Format	Range	Resolution
6	4	[10 4]	$[-2^5, 2^5 - 2^{-4}]$	2^{-4}

- Given the 32-bit floating-point number: **ECE4710A**. Complete the bits in the fields and the significand's FX format: (3 pts.)

sign	e+bias	FX format of significand:
1	110 1100 1	[24 23]
significand		
1 110 0100 0111 0001 0000 1010		

PROBLEM 3 (40 PTS)

- Perform the following 32-bit floating point operations. For fixed-point division, use 4 fractional bits. Truncate the result when required. Show your work: how you got the significand and the biased exponents bits of the result. Provide the 32-bit result.

✓ 4F480000 + D0A90000	✓ C1500000 - C36A0000	✓ FABC0000 × 80400000	✓ 7BB80000 ÷ C9400000
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- ✓ $X = 4F480000 + D0A90000$:

4F480000: 0100 1111 0100 1000 0000 0000 0000 0000

$$e + bias = 10011110 = 158 \rightarrow e = 158 - 127 = 31$$

$$4F480000 = 1.1001 \times 2^{31}$$

$$\text{Significand} = 1.1001$$

D0A90000: 1101 0000 1010 1001 0000 0000 0000 0000

$$e + bias = 10100001 = 161 \rightarrow e = 161 - 127 = 34$$

$$D0A90000 = -1.0101001 \times 2^{34}$$

$$\text{Significand} = 1.0101001$$

$$X = -1.0101001 \times 2^{34} + 1.1001 \times 2^{31} = -1.0101001 \times 2^{34} + \frac{1.1001}{2^3} \times 2^{34}$$

$$X = -(1.0101001 - 0.0011001) \times 2^{34} \text{ (unsigned subtraction)}$$

$$X = -1.001 \times 2^{34}, e + bias = 34 + 127 = 161 = 10100001$$

$$X = 1101 0000 1001 0000 0000 0000 0000 0000 = D0900000$$

b ₈	b ₇	b ₆	b ₅	b ₄	b ₃	b ₂	b ₁	b ₀
1	0	1	0	1	0	0	1	-
0	0	0	1	1	0	0	1	
1	0	0	1	0	0	0	0	

- ✓ $X = C1500000 - C36A0000$:

C1500000: 1100 0001 0101 0000 1000 0000 0000 0000

$$e + bias = 10000010 = 130 \rightarrow e = 130 - 127 = 3$$

$$C1500000 = -1.101 \times 2^3$$

$$\text{Significand} = 1.101$$

C36A0000: 1100 0011 0110 1010 0000 0000 0000 0000

$$e + bias = 10000110 = 134 \rightarrow e = 134 - 127 = 7$$

$$C36A0000 = -1.110101 \times 2^7$$

$$\text{Significand} = 1.110101$$

$$X = -1.101 \times 2^3 + 1.110101 \times 2^7 = -\frac{1.101}{2^4} \times 2^7 + 1.110101 \times 2^7$$

$$X = (-0.0001101 + 1.110101) \times 2^7$$

c ₉	c ₈	c ₇	c ₆	c ₅	c ₄	c ₃	c ₂	c ₁	c ₀
0	1	1	1	0	1	0	1	0	+
1	1	1	1	1	0	0	1	1	
0	1	1	0	1	1	1	0	1	

To subtract these unsigned numbers, we first convert to 2C:

$$R = 01.110101 - 0.0001101 = 01.110101 + 1.1110011$$

The result in 2C is: $R = 01.1011101$

For floating point, we need to convert to sign-and-magnitude:

$$\Rightarrow R(SM) = +1.0111101$$

* You can also do unsigned subtraction: $X = (1.110101 - 0.0001101) \times 2^7$

$$X = 1.1011101 \times 2^7, e + bias = 7 + 127 = 134 = 10000110$$

$$X = 0100 0011 0101 1101 0000 0000 0000 0000 = 435D0000$$

b ₈	b ₇	b ₆	b ₅	b ₄	b ₃	b ₂	b ₁	b ₀
1	1	1	0	1	0	1	0	-
0	0	0	0	1	1	0	1	
1	1	0	1	1	1	0	1	

- ✓ $X = FABC0000 \times 80400000$:

FABC0000: 1111 1010 1011 1100 0000 0000 0000 0000

$$e + bias = 11110101 = 245 \rightarrow e = 245 - 127 = 118$$

$$FABC0000 = -1.01111 \times 2^{118}$$

$$\text{Significand} = 1.01111$$

80400000: 1000 0000 0100 0000 0000 0000 0000 0000

$$e + bias = 00000000 = 0 \rightarrow \text{Denormal number} \rightarrow e = -126 \text{ Significand} = 0.1$$

$$80400000 = -0.1 \times 2^{-126}$$

$$X = (-0.1 \times 2^{-126}) \times (-1.01111 \times 2^{118}) = 0.101111 \times 2^{-8} = 1.01111 \times 2^{-9}$$

$$e + bias = -9 + 127 = 118 = 01110110$$

$$X = 0011 1011 0011 1100 0000 0000 0000 0000 = 3B3C0000$$

- ✓ $X = 7BB80000 \div C9400000$:

7BB80000: 0111 1011 1011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110111 = 247 \rightarrow e = 247 - 127 = 120$$

$$7\text{BB}80000 = 1.0111 \times 2^{120}$$

$$\text{Significand} = 1.0111$$

$$\text{C9400000: } 1100 \ 1001 \ 0100 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000$$

$$e + \text{bias} = 10010010 = 146 \rightarrow e = 146 - 127 = 19$$

$$\text{Significand} = 1.1$$

$$\text{C9400000} = -1.1 \times 2^{19}$$

$$X = -\frac{1.0111 \times 2^{120}}{1.1 \times 2^{19}} = -\frac{1.0111}{1.1} \times 2^{101}$$

$$\begin{array}{r} 000001111 \\ 11000 \overline{) 101110000} \\ \underline{11000} \\ 101100 \\ \underline{11000} \\ 101000 \\ \underline{11000} \\ 100000 \\ \underline{11000} \\ 1000 \end{array}$$

Alignment:

$$\frac{1.0111}{1.1} = \frac{1.0111}{1.1000} = \frac{10111}{11000}$$

Append $x = 4$ zeros: $\frac{101110000}{11000}$

Integer division
 $Q = 1111 \rightarrow Qf = 0.1111$

$$\text{Thus: } X = -0.1111 \times 2^{101} = -1.111 \times 2^{100}$$

$$e + \text{bias} = 100 + 127 = 227 = 11100011$$

$$X = 1111 \ 0001 \ 1111 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 = \text{F1F00000}$$

PROBLEM 4 (30 PTS)

- Fibonacci numbers Computation:** This circuit reads an unsigned number ($n > 1$) and generates F_n :

$$F_n = F_{n-1} + F_{n-2}, F_0 = 0, F_1 = 1$$

- The digital system is depicted below (FSM + Datapath).

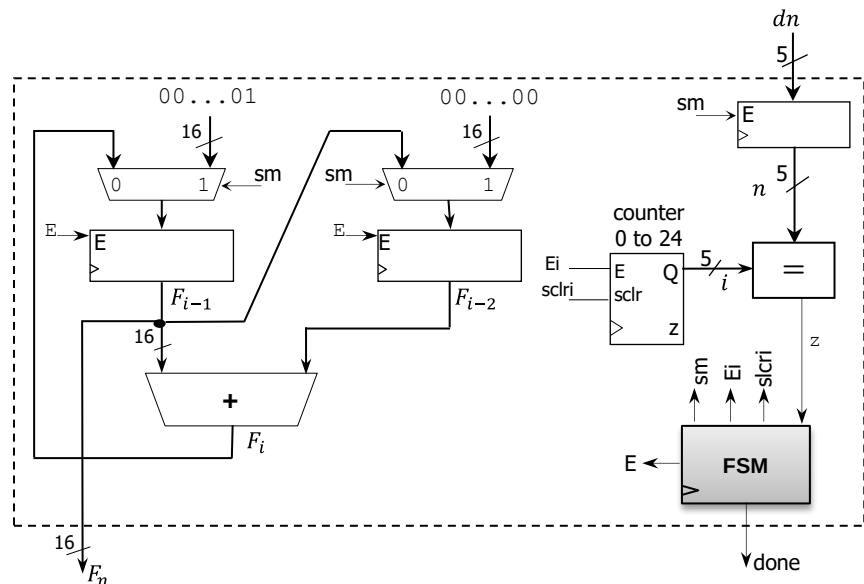
- Input Data dn ($1 < dn < 25$).

Output Data: F_n (16-bit result; $n < 25$ as $F_{24} = 46368$)

- $z = 1$ when $i = n$.

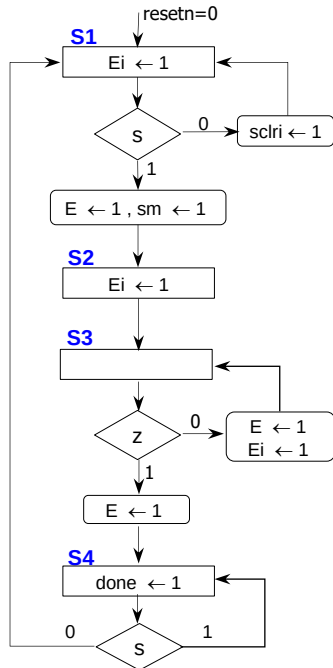
- Sequential Algorithm:

```
n: unsigned integer
F0=0, F1=1
if n > 1
  for i = 2 to n
    Fi = Fi-1 + Fi-2
  end
end
return Fn
```



- Sketch the Finite State Machine diagram (in ASM form) given the sequential algorithm. (15 pts.)
 - The process begins when the s signal is asserted, at this moment we capture dn on register n . At the same time, the registers F_{i-1} and F_{i-2} are initialized with 1 and 0 respectively. Then, the process continues by updating F_{i-1} and F_{i-2} , and it is concluded when $i = n$. The signal $done$ is asserted when the result is ready and appears on output F_n .
 - Note that the iteration index i index must start at $i = 2$.

Finite State Machine:



- Complete the timing diagram. n is provided as an unsigned decimal. You can provide F_{i-1} , F_{i-2} , and i as unsigned decimals. (15 pts.)

